



## Mixed-model Assembly Line Balancing with Reliability

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### ABSTRACT

This paper presents a multi-objective simulated annealing algorithm for the mixed-model assembly line balancing with stochastic processing times. Since, the stochastic task times may have effects on the bottlenecks of a system, maximizing the weighted line efficiency (equivalent to the minimizing the number of station), minimizing the weighted smoothness index and maximizing the system reliability are considered. After solving an example in detail, the performance of the proposed algorithm is examined on a set of test problems. The experimental results show the new approach performs well.

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## 1. INTRODUCTION

An assembly line is a production line that unfinished products move continuously through a sequence of stations that these stations are linked together by a material handling system.

Line balancing is one of the most important aspects of the assembly systems which is defined how tasks should be assigned to the stations subject to precedence constraints.

The first scientific article on the assembly line balancing problem (ALBP) was published by Salveson [1]. Then, many studies have been investigated with different situations, constraints, objective(s) and solving methods. There are several good surveys and taxonomies on the ALBP such as in literatures [2-12].

There are several classifications of ALBP. According to the number of product models that will be assembled on the line, it is divided into single, mixed and multi models.

In the single model, only one type of product, in the mixed-model several models of one type of product and

in multi-model different product types in batches are assembled.

There are two famous objective functions for solving ALBP. One of them is minimization of the number of workstations for the given cycle time (Type-I) and another type is minimization of the cycle time for the given number of workstations (Type-II).

According to the number of objective function(s), we can categorize them to single-objective (i.e., [13] and [14]) and multi-objective (i.e., [15-24]). It is interesting that, recently, multi-objective optimization has attracted the research attention in comparison with single-objective problems [25].

By reviewing the articles that have published for assembly line balancing, it is clear that there are several exact, heuristic and meta heuristic algorithms for solving mixed model assembly line balancing problems. Exact methods can get optimal solution in small-sized problems. Due to the NP-hard class of the ALBP [26], many researchers tried to solve these problems to gain optimum or near optimum solution in reasonable computational time. So, many heuristic and meta heuristic algorithms proposed for ALBP.

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Table 1 shows several articles that used exact, heuristic and meta heuristic algorithms for solving mixed-model assembly line balancing problems.

According to the nature of task times, ALBP is classified into two classes: deterministic and stochastic. Most of the researches in the field of Assembly line balancing assumed that the task times are deterministic [27], but in a realistic manufacturing environment, the task time may be random due to worker fatigue, low skill levels, job dissatisfaction, poorly maintained equipment, defects in raw materials, etc. [28]. Hence, verifying stochastic task time in assembly line balancing will be necessary. There are several papers that investigated stochastic task times for assembly line balancing. For example, Tiacci [29] presented an event and object-oriented simulator for assembly lines. His tool, developed in Java, was capable to simulate mixed model assembly lines, with stochastic task times, parallel stations, fixed scheduling sequences, and buffers within workstations. Also, Cakir et al. [25] proposed an algorithm, based on simulated annealing for multi-objective optimization of a single-model stochastic assembly line balancing problem with parallel stations. The objectives of their paper were (1) minimization of the smoothness index and (2) minimization of the design cost.

A good measure of assembly line balancing in stochastic condition is system reliability. So, there are some papers in this field such as literatures [30-33]. Reliability can be used as a good index when there is uncertainty or probabilistic parameters for system. One of these uncertainties, probabilistic or feasible parameters may be processing times when human involve in assembly line. The reliability of a system with stochastic task time can be defined as a probability that there is no bottleneck in a system.

To the best of our knowledge and literature review, there is no paper that investigated stochastic mixed model assembly line balancing problem according to system reliability, weighted line efficiency and weighted smoothness index, simultaneously. So, this field can be a good area for developing and in this paper we focus on this gap.

**TABLE 1.** Exact, heuristic and meta heuristics for solving mixed-model ALBP

<b>Exact</b>	Branch and Bound	[34]
<b>Heuristics</b>	Heuristic algorithm	[28, 35]
	Simulated Annealing	[16]
	Genetic Algorithm (GA)	[36, 37]
<b>Meta Heuristic</b>	Ant Colony Optimization (ACO)	[17, 18]
	Tabu Search (TS)	[38]
	Particle Swarm Optimization (PSO)	[19]

For this purpose, we propose an SA algorithm for solving mixed model assembly line balancing with stochastic processing time that minimizes weighted smoothness index and maximizes system reliability and weighted line efficiency. The rest of this paper is structured as follows. Section 2 provides some basic concepts about the standard simulated annealing algorithm and weighted sum method for solving multi-objective mathematical models. Problem definition and the proposed simulated annealing algorithm are presented in Section 3. Numerical example and numerical experiments are given in Sections 4 and 5. Finally, Section 6 is devoted to conclusions and recommendations for future research.

## 2. BASIC CONCEPTS

In this section, we introduce the SA algorithm and weighted sum method for solving multi-objective problems.

### 2. 1. The Standard Simulated Annealing Algorithm

The Simulated Annealing algorithm is a random search optimization technique that got its existence from the physical annealing of solid metal.

As Simulated Annealing starts, an initial solution is generated and used as the first current solution. A control parameter ( $T$ ), is specified analogous to the annealing temperature. This temperature is systematically decreased according to a cooling rate. As the temperature drops, neighboring solutions to the current solution are found. If the objective function value is superior to that of the current solution, the neighboring solution becomes the new current solution. If the neighboring solution provides an objective function value inferior to that of the current solution, the neighboring solution may still become the current solution if a certain acceptance criterion is met. A distinctive feature of Simulated Annealing is that inferior solutions are sometimes accepted as the current solution to prevent getting trapped in local optima. Through the occasional acceptance of inferior solutions which meet the acceptance criteria, the search moves to a different location on the continuum of feasible solutions in an effort to reach the global optimum. The process of finding neighboring solutions and accepting these as current solutions if acceptance criteria are met is repeated according to the cooling pattern until some stopping criteria is met [39].

### 2. 2. Weighted Sum Method

This method is one of the most widely used methods for solving multi-objective problems. It composes the set of objectives into a single objective by multiplying each objective with a user supplied weight that this weight depends on the relative importance of each objective. The structure

of this method is given below [40]:

$$\begin{aligned} \text{Min } F(X) &= \sum_{m=1}^M w_m f_m(X) \\ \text{subject to: } G(X) &= [g_1(X), g_2(X), \dots, g_J(X)] \geq 0 \\ H(X) &= [h_1(X), h_2(X), \dots, h_K(X)] = 0 \\ x_i^{(L)} &\leq x_i \leq x_i^{(U)}, i = 1, 2, \dots, N \end{aligned} \quad (1)$$

where, the objectives are normalized and  $w_m \in [0, 1]$  is the weight of the  $m^{\text{th}}$  objective function.

It is usual practice to choose weights such that  $\sum_{i=1}^M W_m = 1$ .

### 3. PROBLEM DEFINITION

In this section the problem assumptions and the proposed algorithm for mixed model assembly line balancing problem with stochastic processing time for maximizing the weighted line efficiency (minimizing the number of stations), minimizing the smoothness index and maximizing the system reliability are introduced.

**3. 1. Problem Assumptions** The assumptions of this problem are given as follows:

1. The required time to do Task  $j$  is stochastic, and it has a Normal distribution with mean  $t_j$  and standard deviation  $\sigma_j$ .
2. Precedence diagrams of different product models are known, and a task cannot be performed until all its predecessors have been completed
3. Common tasks among different product models exist. A task completion time can be different from one model to another.
4. Parallel stations and work-in-process inventories are not allowed.
5. Tasks must be processed only once in each cycle and each task can be assigned to only one station.
6. Stations are arranged in a simple straight assembly line.
7. The maximum cycle time is given.
8. All line workers are paid the same hourly rate and each station is manned by one worker.
9. Demand rate is deterministic.

**3. 2. The Proposed SA Algorithm** In the proposed SA algorithm, the temperature of each iteration is decreased by using the following relation until the final temperature is reached

$$T_{C+1} = \alpha \cdot T_C \quad (2)$$

where,  $\alpha$ ,  $T_C$  and  $T_{C+1}$  are cooling rate, current temperature and next temperature, simultaneously.

Initial solution generation, neighborhood move and structure of building a feasible solution in the algorithm are given as follows.

**3. 2. 1. Initial Solution Generation** Each solution in proposed algorithm is a string of integer numbers.

The initial solution of proposed algorithm is shown in a list that is named priority list ( $PL$ ) and the length of this list is as equal as the number of tasks. The position and the value of the position of this list are important. At the first time, this list generates randomly.

For example if there are 6 tasks in an assembly line, an initial and random priority list can be shown with  $PL = \{2, 1, 4, 5, 3, 6\}$ . It means that Task 2 has the highest priority value and Task 6 has the lowest priority value.

For creating a feasible solution, the assignable tasks that satisfy the precedence constraints are assigned to the station according to their priority values. Then, the set of assignable tasks is updated. Also, when the current station is loaded maximally, it is closed and the next station is opened. This process continues until all tasks are assigned to the stations.

**3. 2. 2. Neighborhood Move** In the proposed algorithm, a neighbor solution of priority list is generated by interchanging 2 or 3 tasks randomly with a probability of 0.5 which is shown in Figure 1. If the generate random value is less than or equal to 0.5, interchanging 2 tasks will be selected, otherwise, interchanging 3 tasks method will be happened.

**3. 2. 3. Building a Feasible Solution** In the procedure of building a feasible solution, the stations have been considered successively. Before the presentation of the procedure of building a feasible solution and calculating the objective functions, it is necessary to introduce the following notations:

- $i, h, p, r$ : Task indices
- $j$ : Station index
- $m$ : Product model
- $M$ : Set of product models
- $P(i)$ : Set of immediate predecessors of Task  $i$
- $t_{im}$ : Operation time of Task  $i$  for model  $m$
- $t_{im}^f$ : Finish time of Task  $i$  for model  $m$
- $NS$ : Number of stations
- $NM$ : Number of models
- $NT$ : Number of tasks
- $SAT$ : Set of assignable tasks
- ${}_m WL_{NS}$ : The station load including unavoidable idle times on the station for all  $m \in M$
- $TL_{NS}$ : The set of tasks which are assigned to the station
- $C$ : Maximum cycle time
- $C_i$ : Trial cycle time
- $C_{min}$ : Minimum cycle time

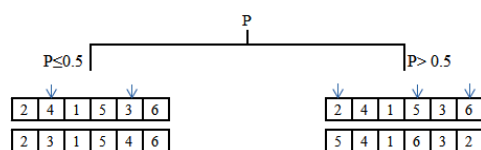


Figure 1. Neighborhood generation

The Procedure of building a solution is as follows:

1. Set  $NS = 1, {}_mWL_{NS}=0$  for all  $m \in M$ .
2. Determine SAT ( $SAT = \{i \mid (\text{all } p \in P(i) \text{ have already been assigned or } P(i) = \{\emptyset\}) \text{ and Task } i \text{ has not been assigned}\}$ ). If  $SAT = \{\emptyset\}$ , then go to Step 6.
3. Sort the tasks in SAT in increasing order of priority value of tasks in PL.
4. Assign the first Task  $h$  in SAT for which;
  - 4.1. If  $t_{hm} + {}_mWL_{NS} \leq C_t$  and  $t_{hm} + t_{rm}^f \leq C_t$  ( $t_{rm}^f = \max \{t_{pml}^f \mid p \in P(h) \text{ have already been assigned to the station}\}$ ) for all  $m \in M$ , then assign Task  $h$  to the station;  $TL_{NS} = TL_{NS} + \{h\}$ , and set  $t_{hm}^f = \max \{(t_{hm} + {}_mWL_{NS}), (t_{hm} + t_{rm}^f)\}$  for all  $m \in M$ . Set  ${}_mWL_{NS} = t_{hm}^f$  for all  $m \in M$  and go to Step 2; otherwise go to Step 5.
5. If none of these tasks in SAT could be assigned at the station, then open a new station. If  $TL_{NS} \neq \{\emptyset\}$  then  $NS = NS + 1, {}_mWL_{NS} = 0$  for all  $m \in M$ , and go to Step 2.
6. Stop.

The trial cycle time ( $C_t$ ) starts from minimum feasible cycle time in the above procedure. It is as follows:  $C_{min} = \max[19 \ i=1, 2, \dots, NT \text{ and } m=1, 2, \dots, NM]$  After creating a feasible solution with this trial cycle time, the objective function according to Section 3.2.4 is calculated. Then the trial cycle time is increased by one unit and the above procedure is repeated until  $C_t \leq C$ .

**3. 2. 4. Objective Function** The objectives of the proposed algorithm for mixed model assembly line balancing with stochastic task time for the given maximum cycle time are as follows:

1. Maximization of the weighted line efficiency.  
It is equivalent to minimize the number of stations or minimizing the line length or the number of operators. Considering the mixed-model nature of the problem, the weighted line efficiency ( $WLE$ ) is calculated as follows for a given line balance [17]:

$$WLE = \left( \frac{\sum_{m \in M} q_m (\sum_{i \in I} t_{im})}{C \cdot NS} \right) \cdot 100 \tag{3}$$

where,  $q_m$  is the overall proportion of the number of units of model  $m$ .  $q_m$  is computed by the following equation where  $D_m$  denotes the demand, over the planning horizon, for model  $m$ .

$$q_m = \frac{D_m}{\sum_{m \in M} D_m} \tag{4}$$

2. Minimizing the weighted smoothness index.  
This index permits decreasing the workload difference between stations where  $WL_{max}$  is the maximum station time.

$$WSI = \sqrt{\frac{\sum_{m \in M} q_m \cdot (\sum_{j \in J} ({}_mWL_j - WL_{max})^2)}{NS}} \tag{5}$$

3. Maximizing the reliability of system.  
In this system, the reliability of each station means the probability that the station is not a bottleneck according to stochastic task time. Thus, reliability of  $j^{th}$

workstation ( $R_j$ ) with trial cycle time  $C_t$  can be defined as follows:

$$R_j = P(\sum_{i=1}^{NT} \sum_{m=1}^{NM} q_m \cdot t_{imj} \leq C_t) = P\left(\frac{(\sum_{i=1}^{NT} \sum_{m=1}^{NM} q_m \cdot t_{imj}) - E(\sum_{i=1}^{NT} \sum_{m=1}^{NM} q_m \cdot t_{imj})}{\sqrt{\sum_{i=1}^{NT} \sum_{m=1}^{NM} q_m^2 \cdot \text{var}(t_{imj})}} \leq \frac{C_t - E(\sum_{i=1}^{NT} \sum_{m=1}^{NM} q_m \cdot t_{imj})}{\sqrt{\sum_{i=1}^{NT} \sum_{m=1}^{NM} q_m^2 \cdot \text{var}(t_{imj})}}\right) = P(Z \leq \frac{C_t - (\sum_{i=1}^{NT} \sum_{m=1}^{NM} q_m \cdot t_{imj})}{\sqrt{\sum_{i=1}^{NT} \sum_{m=1}^{NM} q_m^2 \cdot \text{var}(t_{imj})}}) \tag{6}$$

Since we have an arrangement of  $N$  stations in series, the reliability of the assembly line ( $R_{AL}$ ) can be expressed as:

$$R_{AL} = \prod_{j=1}^N R_j \tag{7}$$

According to the weighted sum method, the objective function of the proposed approach is as follow:

$$\text{Minimize } E = W_1 \left( \frac{WLE_0}{WLE} \right) + W_2 \left( \frac{WSI}{WSI_0} \right) + W_3 \left( \frac{R_{AL_0}}{R_{AL}} \right) \tag{8}$$

where,  $WLE_0$ ,  $WSI_0$  and  $R_{AL_0}$  are the objective function values obtained from the initial solution and  $W_1$ ,  $W_2$  and  $W_3$  are the weights of objectives in the weighted sum method. In this paper, the weight of each objective function is 1.3.

**3. 3. Simple Lower Bound** In this section, we propose a simple lower bound on the minimal number of stations for mixed model stochastic assembly line balancing. This lower bound is as follows:

$$LB = \left\lceil \frac{\sum_{m=1}^{NM} \sum_{i=1}^{NT} q_m \cdot t_{im}}{C_t} \right\rceil \tag{9}$$

( $\lceil x \rceil$  denotes the smallest integer not being smaller than  $x$ ).

**3. 4. Parameter Settings** In the meta heuristic algorithms, choosing the best combination of the parameters can intensify the search process and prevent premature convergence.

In this paper, the Taguchi (1986) method is used for the best parameter selections.

Three levels are selected for each parameter of the SA algorithm. They are shown in Table 2.

The Taguchi method uses orthogonal arrays for decreasing the number of experiments for parameter settings. These arrays are presented in Table 3.

**TABLE 2.** Factors and their levels

Factor	Initial temperature			Final temperature			Length of the Markov chain			Cooling rate		
level	1	2	3	1	2	3	1	2	3	1	2	3
value	50	100	150	0.5	1	2	5	10	n*	0.9	0.95	0.99

n\*: Number of tasks

**TABLE 3.** The orthogonal arrays for the proposed approach

Test	Initial Temperature	Final Temperature	Length Of the Markov Chain	Cooling Rate
1	1	1	1	1
2	1	2	2	2
3	1	3	3	3
4	2	1	2	3
5	2	2	3	1
6	2	3	1	2
7	3	1	3	2
8	3	2	1	3
9	3	3	2	1

It shows nine tests are necessary to select the best value for each parameter.

Each test is run four times, and the average of the objective function is obtained to estimate the (SN) ratio. In the Taguchi method, the S/N ratio is as follows:

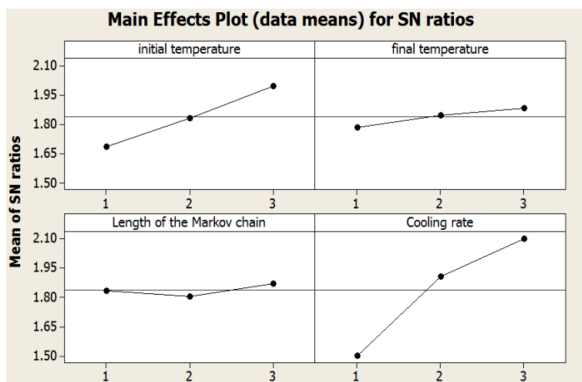
$$SN = -10\log\left(\frac{1}{n}\sum_{i=1}^n(\text{objective function})^2\right) \quad (10)$$

Each level which has the maximum SN ratio is the best one.

According to Figure 2, the best level of each parameter is reported in Table 4.

**4. NUMERICAL EXAMPLE**

We illustrate the proposed algorithm by using a nine-task and two-model example problem. Expected task times and their variances are generated randomly.



**Figure 2.** The mean SN ratio plot for the selected levels of each factor

**TABLE 4.** Factors and their levels

Factor	Initial Temperature	Final Temperature	Length of the Markov Chain	Cooling Rate
level	3	3	3	3
value	150	2	n*	0.99

The required data such as Expected task time ( $\mu(i)$ ) and variance of task time ( $\sigma_i^2$ ) of this example are given in Table 5. The maximum cycle time of this problem is 9.

The overall proportion of the number of units of model A and B is 0.5. So,  $q_A=q_B=50\%$ . The initial random solution (priority list) constructed as:  $PL = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$ . The procedure of creating the initial line balance is shown in Table 6. The assignment of tasks to the stations and the reliability of each station are presented in Table 7. It shows there are seven stations in system with initial trial cycle time=3. The reliability of station 5 is lower than the others. It can show the importance of this station because it has this ability that will be a bottleneck. The objective function values of WLE, WSI, RLA and E of the initial line balance are 59.524%, 1.711, 0.389 and 1, respectively.

In the next step, a new neighbor solution is generated by interchanging 2 or 3 tasks randomly with a probability of 0.5. These steps are repeated until the final temperature is met. Then, the trial cycle time is increased by one unit and the above procedure is repeated until  $C_t \leq 9$ .

In this problem, according to several preliminary experiments we selected initial temperature, final temperature and cooling rate as 100, 1 and 0.95, respectively.

We run this algorithm 5 times with PC 2.2 GHz CPU and 1 GB of RAM. The best and the average results of these iterations are presented in Table 8.

The best function value in 5 iterations with different initial random solution for this problem is 0.509. The number of stations is 5 and the RLA, WSI and WLE are 0.100, 0.949 and 83.333, respectively.

**5. NUMERICAL EXPERIMENT**

In order to assess the effectiveness of the proposed algorithm, a set of standard test problems (P9, P14, P20, P25, P30, P39, P47 and P65) are solved.

**TABLE 5.** Data of the example problem

Task	Immediate Predecessors	Model A		Model B	
		$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$
1	—	2	0.5	0	0
2	—	3	0.8	1	0.3
3	—	0	0	1	0.3
4	1	3	0.8	0	0
5	2	1	0.3	3	0.8
6	2,3	1	0.3	1	0.3
7	4,5	2	0.5	2	0.5
8	5	0	0	3	0.8
9	6	1	0.3	1	0.3

**TABLE 6.** Building the initial line balance

Step1	Step2	Step3	Step4	Step5	Step6
NS=1; ${}^A WL_1=0$ ; ${}^B WL_1=0$ .	SAT={1,2,3}	PL={1,2,3}	Select Task 1, $P(1) = \{\emptyset\}$ ; $2+0 \leq 3$ ; $0+0 \leq 3$ ; $TL_1 = TL_1 + \{1\}$ ; $t_{1A}^f=2, t_{1B}^f=0, {}^A WL_1=2, {}^B WL_1=0$		
	SAT={2,3,4}	PL={2,3,4}	Select Task 2, $P(2) = \{\emptyset\}$ ; $2+3 > 3$ ; go to step 5	Task2 could not be selected	
NS=2; ${}^A WL_2=0$ ; ${}^B WL_2=0$	SAT={2,3,4}	PL={2,3,4}	Select Task 2, $P(2) = \{\emptyset\}$ ; $3+0 \leq 3$ ; $1+0 \leq 3$ ; $TL_2 = TL_2 + \{2\}$ ; $t_{2A}^f=3, t_{2B}^f=1, {}^A WL_2=3, {}^B WL_2=1$		
	SAT={3,4,5}	PL={3,4,5}	Select Task 3, $P(3) = \{\emptyset\}$ ; $0+3 \leq 3$ ; $1+1 \leq 3$ ; $TL_2 = TL_2 + \{3\}$ ; $t_{3A}^f=3, t_{3B}^f=2, {}^A WL_2=3, {}^B WL_2=2$		
	SAT={4,5,6}	PL={4,5,6}	Select Task 4, $P(4) = \{1\}$ ; $3+3 > 3$ ; go to step 5	Task4 could not be selected	
NS=3; ${}^A WL_3=0$ ; ${}^B WL_3=0$ .	SAT={4,5,6}	PL={4,5,6}	Select Task 4, $P(4) = \{1\}$ ; $0+3 \leq 3$ ; $0+0 \leq 3$ ; $TL_3 = TL_3 + \{4\}$ ; $t_{4A}^f=3, t_{4B}^f=0, {}^A WL_3=3, {}^B WL_3=0$		
	SAT={5,6}	PL={5,6}	Select task 5, $P(5) = \{2\}$ ; $1+3 > 3$ ; go to step 5	Task5 could not be selected	
NS=4; ${}^A WL_4=0$ ; ${}^B WL_4=0$	SAT={5,6}	PL={5,6}	Select Task 5, $P(5) = \{2\}$ ; $1+0 \leq 3$ ; $3+0 \leq 3$ ; $TL_4 = TL_4 + \{5\}$ ; $t_{5A}^f=1, t_{5B}^f=3, {}^A WL_4=1, {}^B WL_4=3$		
	SAT={6,7,8}	PL={6,7,8}	Select Task 6, $P(6) = \{2,3\}$ ; $3+3 > 3$ ; go to step 5	Task6 could not be selected	
NS=5; ${}^A WL_5=0$ ; ${}^B WL_5=0$ .	SAT={6,7,8}	PL={6,7,8}	Select Task 6, $P(6) = \{2,3\}$ ; $1+0 \leq 3$ ; $1+0 \leq 3$ ; $TL_5 = TL_5 + \{6\}$ ; $t_{6A}^f=1, t_{6B}^f=1, {}^A WL_5=1, {}^B WL_5=1$		
	SAT={7,8,9}	PL={7,8,9}	Select Task 7, $P(6) = \{4,5\}$ ; $2+1 \leq 3$ ; $2+1 \leq 3$ ; $TL_5 = TL_5 + \{7\}$ ; $t_{7A}^f=3, t_{7B}^f=3, {}^A WL_5=3, {}^B WL_5=3$		
NS=6; ${}^A WL_6=0$ ; ${}^B WL_6=0$ .	SAT={8,9}	PL={8,9}	Select Task 8, $P(8) = \{5\}$ ; $3+3 > 3$ ; go to step 5	Task8 could not be selected	
	SAT={8,9}	PL={8,9}	Select Task 8, $P(8) = \{5\}$ ; $0+0 \leq 3$ ; $3+0 \leq 3$ ; $TL_6 = TL_6 + \{8\}$ ; $t_{8A}^f=0, t_{8B}^f=3, {}^A WL_6=0, {}^B WL_6=3$		
	SAT={9}	PL={9}	Select Task 9, $P(9) = \{6\}$ ; $1+3 > 3$ ; go to step 5	Task9 could not be selected	
NS=7; ${}^A WL_7=0$ ; ${}^B WL_7=0$ .	SAT={9}	PL={9}	Select Task 9, $P(9) = \{6\}$ ; $1+0 \leq 3$ ; $1+0 \leq 3$ ; $TL_7 = TL_7 + \{9\}$ ; $t_{9A}^f=1, t_{9B}^f=1, {}^A WL_7=1, {}^B WL_7=1$		
	SAT={ $\emptyset$ }				Stop

**TABLE 7.** The reliability of each station

Station	1	2	3	4	5	6	7
Tasks	1	2,3	4	5	6,7	8	9
$\sum_{i=1}^{NT} \sum_{m=1}^{NM} q_m \mu_{imj}$	1	2.5	1.5	2	3	1.5	1
$\sum_{i=1}^{NT} \sum_{m=1}^{NM} q_m^2 var(t_{imj})$	0.125	0.35	0.2	0.275	0.4	0.2	0.15
$RL_j$	1.0000	0.801	0.9996	0.9717	0.5	0.9996	1.0000

The details of these problems are illustrated in Appendix. The parameters of the proposed algorithm are as follows:

$T_0=100$ ;  $T_0=1$ ;  $r=0.95$  and the length of Markov chain is as equal as the number of tasks. Each problem is solved five times with initial random solution and the

best and average solutions for each trial cycle times are reported. Also, the lower bounds of the number of stations for each trial cycle times are calculated. These results are shown at Tables 8 and 9 and it is possible to compare the performance of the proposed algorithm with the LB.

**TABLE 8.** Comparison results for the small-sized test problems.

Trial cycle time	LB	Best					Average				Elapsed Time(s)
		E	NS	RLA	WSI	WLE	E	RLA	WSI	WLE	
<b>3</b>	5	<b>0.509</b>	<b>5</b>	<b>0.100</b>	<b>0.949</b>	<b>83.333</b>	0.554	0.154	0.949	83.333	
<b>4</b>	4	0.715	4	0.249	1.275	78.125	0.718	0.243	1.313	78.125	
<b>5</b>	3	1.000	4	0.389	1.711	59.524	1.000	0.389	1.711	59.524	
<b>P9</b>	<b>6</b>	1.000	4	0.389	1.711	59.524	1.000	0.389	1.711	59.524	2.78
	<b>7</b>	0.851	2	0.479	1.118	89.286	0.851	0.479	1.118	89.286	
	<b>8</b>	1.000	2	0.389	1.711	59.524	1.000	0.389	1.711	59.524	
	<b>9</b>	1.000	2	0.389	1.711	59.524	1.000	0.389	1.711	59.524	
	<b>8</b>	0.749	10	0.364	3.375	66.700	0.831	0.357	3.763	61.849	
	<b>9</b>	0.674	8	0.343	3.311	74.111	0.744	0.422	3.345	68.182	
	<b>10</b>	<b>0.616</b>	<b>7</b>	<b>0.538</b>	<b>3.104</b>	<b>76.229</b>	0.695	0.471	3.387	76.229	
	<b>11</b>	0.677	7	0.705	3.633	69.299	0.685	0.775	3.528	69.299	
<b>P14</b>	<b>12</b>	0.630	6	0.864	3.485	74.111	0.670	0.788	3.604	74.111	10.88
	<b>13</b>	0.649	6	0.958	3.516	68.410	0.698	0.875	3.713	68.410	
	<b>14</b>	0.684	5	0.669	4.020	76.229	0.712	0.668	4.020	76.229	
	<b>15</b>	0.702	5	0.836	4.190	71.147	0.717	0.850	4.054	71.147	
	<b>16</b>	0.698	5	0.949	4.020	66.700	0.740	0.959	4.174	66.700	
	<b>12</b>	0.816	10	0.638	3.476	76.000	0.892	0.690	3.596	74.618	
	<b>13</b>	<b>0.724</b>	<b>9</b>	<b>0.811</b>	<b>3.242</b>	<b>77.949</b>	0.810	0.835	3.289	76.390	
<b>P20</b>	<b>14</b>	0.730	9	0.925	3.525	72.381	0.831	0.942	3.752	72.381	21.03
	<b>15</b>	0.760	8	0.886	4.099	76.000	0.879	0.913	4.185	70.933	
	<b>16</b>	0.844	9	0.929	4.752	63.333	0.933	0.910	4.657	66.500	
	<b>17</b>	0.877	8	0.911	5.532	67.059	0.984	0.914	5.406	67.059	
	<b>20</b>	0.904	18	0.999	9.741	58.278	0.921	0.999	9.783	57.664	
	<b>21</b>	0.918	17	0.900	9.430	58.768	0.933	0.949	9.701	57.462	
	<b>22</b>	0.886	15	0.962	9.656	63.576	0.936	0.971	10.448	60.397	
<b>P25</b>	<b>23</b>	0.908	14	0.888	10.009	65.155	0.949	0.941	10.529	60.160	37.87
	<b>24</b>	<b>0.884</b>	<b>13</b>	<b>0.914</b>	<b>9.717</b>	<b>67.244</b>	0.929	0.949	10.248	61.736	
	<b>25</b>	0.925	13	0.927	10.693	64.554	0.953	0.974	10.787	59.267	
	<b>26</b>	0.924	14	0.998	10.227	57.637	0.944	0.989	10.546	58.642	
	<b>27</b>	0.935	14	0.998	10.218	55.503	0.979	0.986	11.202	56.357	

The above table shown the proposed algorithm can be as an effective algorithm because the initial objective value (E) was 1 and it decreased the duration of running algorithm. Furthermore, the weighted line efficiency and the reliability of system were increased and the weighted smoothness index was decreased, simultaneously. For example, the initial values of WLE, WSI and RLA for each P65 are given as follows: Also, the number of stations found by SA algorithm is compared to LB given in Equation (10).

As it can be seen, the proposed SA algorithm performs well throughout on the different problems. Figure 3 shows a comparison between the lower bound and the obtained number of stations by the proposed algorithm. This Figure shows the structure of the problem (predecessors, task times, ...) has important effect on the obtained results.

So, there is no regular procedure for number of stations and cycle time. However, by increasing LB, NS increases, too.

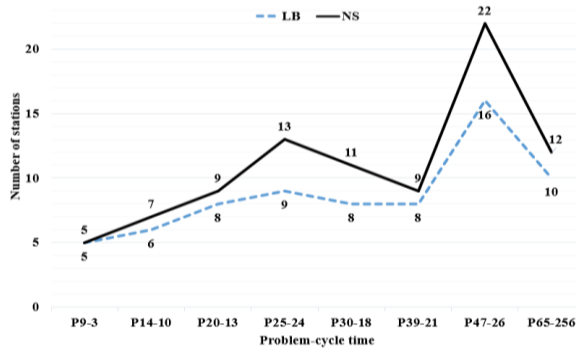
TABLE 9. Computational results for the large-sized test problems

Trial cycle time	LB	Best					Average				Elapsed Time(s)
		E	NS	RLA	WSI	WLE	E	RLA	WSI	WLE	
12	12	0.667	16	0.293	3.903	72.969	0.745	0.281	4.130	71.252	50.02
13	11	0.656	15	0.384	3.603	71.846	0.705	0.404	4.279	70.948	
14	11	0.631	14	0.472	3.951	71.480	0.679	0.416	4.031	72.579	
15	10	0.619	13	0.639	4.126	71.846	0.664	0.567	4.187	73.044	
16	9	0.644	11	0.264	3.701	79.602	0.675	0.324	3.858	78.276	
17	9	0.621	11	0.476	4.023	74.920	0.680	0.547	4.452	73.671	
18	8	<b>0.582</b>	<b>11</b>	<b>0.866</b>	<b>3.704</b>	<b>70.758</b>	0.633	0.773	4.132	73.588	
19	8	0.604	10	0.836	4.226	73.737	0.659	0.771	4.568	75.375	
20	8	0.635	10	0.967	4.557	70.050	0.653	0.719	4.432	74.720	
13	13	0.657	16	0.113	3.978	76.779	0.736	0.168	4.111	75.093	
14	12	0.600	15	0.311	3.414	76.048	0.680	0.341	3.985	73.308	
15	11	0.686	14	0.142	4.580	76.048	0.750	0.217	4.742	75.034	
16	10	0.671	13	0.192	4.634	76.779	0.759	0.352	4.921	71.441	
17	10	0.660	12	0.213	4.619	78.284	0.739	0.371	4.669	75.234	
18	9	0.635	11	0.169	4.242	80.657	0.715	0.208	4.526	80.657	
19	9	0.603	11	0.604	3.574	76.412	0.719	0.502	4.810	76.412	
20	8	0.645	10	0.481	4.797	79.850	0.723	0.432	5.001	79.850	
21	8	<b>0.598</b>	<b>9</b>	<b>0.324</b>	<b>4.236</b>	<b>84.497</b>	0.692	0.256	4.455	84.497	
22	8	0.677	9	0.428	5.279	80.657	0.759	0.375	5.449	80.657	
23	18	0.801	26	0.535	9.514	67.475	0.867	0.568	9.756	65.511	158.67
24	17	0.716	25	0.786	8.702	67.250	0.785	0.810	9.225	65.181	
25	17	0.722	24	0.861	9.219	67.250	0.760	0.904	9.426	68.420	
26	16	<b>0.685</b>	<b>22</b>	<b>0.968</b>	<b>8.921</b>	<b>70.542</b>	0.793	0.857	9.695	65.578	
27	15	0.718	22	0.891	9.325	67.929	0.782	0.843	9.836	69.223	
28	15	0.728	21	0.937	9.675	68.622	0.790	0.868	10.140	68.685	
29	14	0.731	20	0.985	10.309	69.569	0.778	0.935	10.321	69.569	
30	14	0.750	19	0.776	10.202	70.790	0.785	0.880	10.259	69.374	
31	14	0.753	18	0.821	10.704	72.312	0.797	0.907	10.786	70.104	
32	13	0.742	18	0.942	10.556	70.052	0.787	0.936	10.644	70.139	
33	13	0.737	17	0.919	10.579	71.925	0.788	0.910	10.731	71.126	448.76
34	12	0.739	16	0.985	11.133	74.173	0.802	0.909	11.295	72.428	
35	12	0.724	15	0.857	10.481	76.857	0.785	0.905	10.969	73.975	
249	10	0.896	13	0.999	74.316	75.779	0.960	0.999	74.954	75.779	
250	10	0.901	13	0.999	75.561	75.475	0.966	0.999	76.172	75.475	
251	10	0.927	13	1.000	82.768	75.175	0.968	1.000	76.694	75.175	
252	10	0.920	13	1.000	80.392	74.876	0.971	1.000	77.075	74.876	
253	10	0.922	13	0.999	80.521	74.580	0.983	1.000	79.469	74.580	
254	10	0.906	13	1.000	75.478	74.287	0.978	1.000	77.852	74.287	
255	10	0.911	13	1.000	76.789	73.996	0.981	0.999	78.164	73.996	
256	10	<b>0.855</b>	<b>12</b>	<b>0.999</b>	<b>66.925</b>	<b>79.849</b>	0.936	0.978	69.899	78.620	
257	10	0.913	13	1.000	76.474	73.420	0.976	0.991	76.864	74.643	
258	10	0.908	13	1.000	74.641	73.135	0.984	0.988	76.976	73.135	
259	10	0.943	12	0.981	71.982	78.924	0.990	0.956	77.727	75.281	
260	10	0.885	12	0.958	70.250	78.620	0.980	0.987	77.727	74.992	



**TABLE 10.** Comparison between the initial and the best objective functions for P65

WL E0	RL A0	W SI0	W LE	RL A	W SI	WLE/ WLE0	RLA/ RLA0	WSI/ WSI0
70.	0.9	97.	79.	0.9	66.			
36	98	94	84	99	92	↑	↑	↓

**Figure 3.** Comparison between the lower bound and the obtained number of stations

## 6. CONCLUSION

In this paper, we presented a multi-objective simulated annealing algorithm for mixed-model assembly line balancing with stochastic processing time to maximize the weighted line efficiency (minimizing number of stations), minimizing the weighted smoothness index and maximizing the reliability of system. In this problem maximum cycle time is given. An illustrative example problem is solved by using the proposed algorithm, and numerical experiments are conducted to demonstrate the efficiency of the proposed approach. The results show that the proposed approach obtains good solutions within a short computational time for every test problem because the best result of the objective value (E) in the initial solution was 1 and it decreased in the duration of the proposed algorithm. For further researches the development of this condition for a given number of stations and also using the other meta heuristics may be good subjects.

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**8. APPENDIX**

The details of task times and variances for each problem are presented in Tables A1, A2, A3, A4, A5, A6 and A7.

**TABLE A1. Problem P14**

Task	Immediate predecessor(s)	Model A $q_A=0.42$		Model B $q_B=0.58$		Task	Immediate predecessor(s)	Model A $q_A=0.42$		Model B $q_B=0.58$	
		$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$			$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$
1	---	0	0	2	0.5	8	4,5	0	0	2	0.5
2	---	8	4	8	4	9	5	3	1.5	2	0.5
3	1	7	3	7	3	10	6	3	1.5	2	0.5
4	3	7	3	5	2	11	5,6	6	2.5	6	2.5
5	3	2	0.5	2	0.5	12	8	3	1.5	3	1.5
6	3	6	2.5	0	0	13	7, 10, 11	5	2	5	2
7	2,3	4	2	0	0	14	9, 12, 13	4	2	6	2.5

**TABLE A2. Problem P20**

Task	Immediate predecessor(s)	Model A $q_A=0.4$		Model B $q_B=0.6$		Task	Immediate predecessor(s)	Model A $q_A=0.4$		Model B $q_B=0.6$	
		$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$			$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$
1	---	4	1	0	0	11	8	3	0.7	4	1
2	---	5	2	4	1	12	9	8	1	7	4
3	---	0	0	2	0.6	13	10	2	0.2	0	0
4	1	2	1	0	0	14	11	11	4	10	2
5	1,2	4	0.3	5	1	15	11,12	6	1	7	1
6	2,3	3	0.1	4	0.8	16	13,14	10	5	9	5
7	4	5	2	0	0	17	15	12	2	0	0
8	5	5	1	5	1	18	15	0	0	9	2
9	6	7	0.5	8	0.7	19	16,17,18	4	1	4	1
10	7	6	2	0	0	20	19	5	1	6	3

**TABLE A3. Problem P25**

Task	Immediate predecessor(s)	Model A $q_A=0.4$		Model B $q_B=0.6$		Task	Immediate predecessor(s)	Model A $q_A=0.4$		Model B $q_B=0.6$	
		$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$			$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$
1	—	0	0	18	3	14	9	7	1	0	0
2	—	10	4	19	2	15	12,13	17	3	14	2
3	1	15	5	10	1	16	10,13	18	5	11	1
4	3	12	3	10	1	17	16	4	1	0	0
5	3	8	2	3	2	18	16	9	2	6	2
6	3	9	7	0	0	19	14,18	10	3	5	1
7	3	20	5	0	0	20	7,18	0	0	9	2
8	4,5	0	0	2	0.5	21	17	12	2	2	0.3
9	5	15	3	9	1	22	21	18	4	11	1
10	2,6	7	2	12	3	23	15,19,21	12	1	5	1
11	5,6	4	1	10	4	24	20,22,23	10	1	9	1
12	8,9	11	2	10	4	25	24	7	2	0	0
13	11	9	2	12	2						

**TABLE A4. Problem P30**

Task	Immediate predecessor(s)	Model A $q_A=0.5$		Model B $q_B=0.5$		Task	Immediate predecessor(s)	Model A $q_A=0.5$		Model B $q_B=0.5$	
		$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$			$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$
1	—	9.5	3.5	9.5	3.5	16	3	1.4	0.5	1.4	0.5
2	—	1.3	0.5	1.3	0.5	17	3	7.8	3.5	7.8	3
3	—	4.8	2	4.8	2	18	17	2.9	1	2.9	1
4	1	3.3	2	3.3	2	19	18	1.6	0.5	1.6	0.5
5	1	1.5	0.5	1.7	0.5	20	14,16	7	3	7	3
6	5	4.5	2	4.1	2	21	20	8.7	4	8.7	4
7	4, 6	3.6	2	3.6	2	22	15,21	3.9	2	4.1	2
8	7	0	0	2	1	23	22	6.4	3	6.4	3
9	8	12	5	12	5	24	10,20	2.8	1	2.7	1
10	—	0	0	8	3	25	24	8.5	3	8.5	3
11	2	2.5	1	2.5	1.5	26	9,25	6.7	3	6.7	3
12	2	4.3	2	4.3	2	27	23,26	1.9	1	1.9	0.5
13	12	6.5	3	0	0	28	27	9.9	4	9.9	4
14	13	1.7	0.5	1.7	0.5	29	27	4.6	2.2	0	0
15	14	7	3	7	3	30	29	4	2	4.2	2

**TABLE A5. Problem P39**

Task	Immediate predecessor(s)	Model A $q_A=0.45$		Model B $q_B=0.55$		Task	Immediate predecessor(s)	Model A $q_A=0.45$		Model B $q_B=0.55$	
		$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$			$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$
1	—	2	0.6	2	0.8	21	5,18	3	1	3	1.5
2	1	2	0.4	2	1	22	20	8	2.5	8	2.5
3	—	2	0.4	2	0.9	23	—	5	1	5	2
4	—	2	0.3	2	0.5	24	22	7	3	7	3.5
5	—	2	0.4	2	0.3	25	24	4	1	4	2.4
6	2	0	0	11	4	26	25	6	3.4	6	3.4
7	2	0	0	0	0	27	26,23,21	5	2.5	5	1.8
8	6,7	9	2.8	12	3	28	25	0	0	0	0
9	—	2	1	2	1	29	27	1	0.5	1	0.7
10	3,9	10	3	10	3	30	28,29	3	1	3	3.2
11	3	3	1	0	0	31	30	3	1.8	3	1.5
12	8	11	2.5	11	2.3	32	31	0	0	0	0
13	3	4	1.5	4	1.4	33	24	4	1	4	1
14	—	0	0	4	2.3	34	22	2	0.3	2	0.6
15	2	9	3	9	3.5	35	32,33,34	2	0.5	2	0.5
16	15,14,13	13	4	13	3.8	36	35	1	0.7	1	0.4
17	4,11,16	6	2	6	1.9	37	34	1	0.5	1	0.4
18	17	7	3	7	3.5	38	36,37	1	0.9	1	0.6
19	—	3	1.2	3	1.4	39	38	1	0.2	1	0.3
20	10,19	8	1.5	7	2.5						

**TABLE A6. Problem P47**

Task	Immediate predecessor(s)	Model A $q_A=0.45$		Model B $q_B=0.55$		Task	Immediate predecessor(s)	Model A $q_A=0.45$		Model B $q_B=0.55$	
		$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$			$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$
1	—	23	2	2	0.6	25	24	6	2	11	2
2	6	20	3	5	1	26	25	9	1	7	3
3	6	2	0.2	9	2	27	6	7	3	8	1
4	6	6	1	7	0.3	28	20,21	4	1	14	4
5	—	14	2	11	4	29	23	3	0.4	19	6
6	1	22	1	23	5	30	28	2	0.1	11	1
7	6	1	0.1	2	0.1	31	23,27	12	2	15	0.5
8	5,6	7	1	6	2	32	31	13	3	4	0.1
9	6	4	2	8	4	33	34	1	0.4	3	1
10	12	8	1	7	1	34	—	5	2	2	0.4
11	6	12	4	11	3	35	34	4	3	8	0.2
12	16	9	2	14	5	36	6,33,35	13	6	6	1
13	16	7	1	18	6	37	7	18	5	19	0.5
14	16	3	1	3	0.1	38	37	20	8	15	1
15	6	11	3	1	0.3	39	6	8	0.8	3	0.4
16	15	20	6	6	2	40	7,41	11	5	7	0.3
17	7,9	2	0.8	4	1	41	—	17	8	1	0.1
18	17	9	3	5	0.5	42	—	3	0.4	8	3
19	6	7	1	11	3	43	7,36	9	5	9	2
20	17	4	1	9	1	44	36,42	7	3	10	2
21	17	3	0.5	4	0.6	45	44	17	3	20	5
22	17	7	0.2	6	2	46	45	14	4	3	0.3
23	28	11	3	2	0.4	47	44	11	5	3	0.1
24	23	5	2	1	0.1						

**TABLE A7. Problem P65**

Task	Immediate predecessor(s)	Model A $q_A=0.45$		Model B $q_B=0.55$		Task	Immediate predecessor(s)	Model A $q_A=0.45$		Model B $q_B=0.55$	
		$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$			$\mu(i)$	$\sigma_i^2$	$\mu(i)$	$\sigma_i^2$
1	—	35	6	26	3	34	7,35	41	7	39	12
2	3	14	2	15	4	35	23	21	5	63	7
3	4	1	0.2	54	8	36	8	33	7	28	3
4	5	18	3	1	0.2	37	24,36	147	30	103	20
5	6	36	5	33	5	38	10,37	43	12	12	3
6	7	29	4	26	9	39	11,38	15	6	24	12
7	1	159	10	61	12	40	25,39	27	9	1	0.1
8	1	70	8	5	1	41	13,40	4	1	5	0.4
9	8	24	5	16	3	42	26,41	25	5	80	5
10	9	99	9	40	7	43	15,42	22	4	18	9
11	10	56	5	56	7	44	27,43	3	0.4	7	3
12	11	51	4	47	3	45	17,44	44	1	35	8
13	12	94	10	132	18	46	18,28	26	5	34	5
14	13	29	9	6	1	47	30	2	0.1	19	4
15	14	39	12	1	0.2	48	31	41	8	44	12
16	15	15	8	16	6	49	46	9	3	3	0.4
17	16	11	3	2	0.3	50	49	8	2	46	3
18	17	74	15	58	2	51	28,50	20	5	13	2
19	18	34	12	14	8	52	47,62	35	7	13	4
20	21	19	2	19	9	53	48,63	14	3	5	0.8
21	31	14	5	14	6	54	33	38	7	93	9
22	6	12	1	24	3	55	35	11	3	23	3
23	1	19	5	61	7	56	37	56	9	7	1
24	9	38	7	76	12	57	40,64	23	3	98	8
25	12	89	9	43	11	58	42	52	8	64	4
26	14	66	8	45	5	59	44	16	1	2	0.6
27	16	27	3	6	1	60	50	10	0.3	24	3
28	19	189	17	249	25	61	29,62	98	31	66	12
29	20,30	2	0.3	48	7	62	—	3	0.1	11	9
30	21,31	21	6	33	5	63	—	117	25	15	5
31	2,4	6	0.8	5	1	64	—	54	7	85	4
32	5,31,33	18	3	12	7	65	62	35	4	26	3
33	22,34	13	8	53	4						

## Mixed-model Assembly Line Balancing with Reliability

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این مقاله به ارائه یک الگوریتم انجاماد تدریجی چندهدفه برای بالانس خطوط مونتاژ مدل‌های ترکیبی با زمان‌های پردازش تصادفی می‌پردازد. از آنجا که زمان‌های تصادفی ممکن است روی گلوگاه‌های سیستم تأثیرگذار باشد، حداقل سازی تعداد ایستگاه‌ها (معادل با ماکزیمم سازی کارایی موزون خط، حداقل سازی شاخص هموارسازی موزون و ماکزیمم سازی قابلیت اطمینان می‌باشد در این تحقیق مورد مطالعه قرار گرفته است. پس از حل یک مسئله با جزئیات، عملکرد الگوریتم به کمک مجموعه‌ای از مسائل مورد ارزیابی قرار گرفته است که نتایج آزمایشات نشان دهنده عملکرد خوب الگوریتم پیشنهادی است.

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